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Transmitted herewith for filing is the patent application of:

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FOR: SYSTEM AND METHOD FOR CONTROLLING PROGRAM EXECUTION

Enclosed are:

- ☒ 30 pages of specification, claims, abstract.
- ☒ Declaration and Power of Attorney.
- ☒ Priority Claimed.
- ☒ Certified copy of Japanese Patent Application No. 11-029568
- ☒ 12 sheets of formal drawing.
- ☒ An assignment of the invention to Matsushita Electric Industrial Co., Ltd.
and the assignment recordation fee.
- ☒ An associate power of attorney.
- ☐ A verified statement to establish small entity status under 37 CFR 1.9 and 37 CFR 1.27.
- ☒ Information Disclosure Statement, Form PTO-1449 and reference.
- ☒ Return Receipt Postcard
- ☐

The filing fee has been calculated as shown below:

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Respectfully submitted,

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SYSTEM AND METHOD FOR CONTROLLING PROGRAM EXECUTION

BACKGROUND OF THE INVENTION

The present invention relates to controlling the
5 execution of a program in an information processor for
processing instructions by pipelining.

In an information processor such as microprocessor or
digital signal processor (DSP) for executing instructions by
pipeline processing, a pipeline hazard resulting from
10 conditional branching, i.e., a branch hazard, will happen.
Particularly when there is a great number of pipeline stages,
the branch hazard causes a serious problem.

To minimize the issuance of conditional branch
instructions, conditionally executable instructions are
15 adoptable effectively. For example, according to the
technique disclosed in Japanese Laid-Open Publication No. 10-
49368, each conditionally executable instruction contains a
condition field specifying its own execution condition. That
is to say, the instruction is selectively executed depending
20 on whether or not the execution condition specified by its
own condition field is satisfied. This technique is hard to
apply to an information processor of the type using a short
instruction word length, because each instruction to be
executed by the processor cannot afford including such an
25 additional condition field.

As for the short-word-length information processor, an instruction issued exclusively to control the conditional execution of succeeding instructions, i.e., an execution control instruction, is applicable effectively. For instance, according to the technique disclosed in Japanese Laid-Open Publication No. 7-253882, the execution control instruction contains a condition field specifying multiple registers. The number of registers specified in the condition field is always equal to the number of instructions, which succeed the execution control instruction and will be executed under controlled conditions, and is a fixed number. And it is determined based on the value of each of these registers whether or not a succeeding instruction corresponding to the register should be executed. As a result, conditional branch instructions do not have to be used so frequently.

According to the technique disclosed in the latter publication, however, if the number of instructions to be executed under controlled conditions should be increased, then the number of registers specified in the condition field of the execution control instruction has to be increased. As a result, the length of each instruction word must be increased unintentionally. Thus, the short-word-length information processor cannot control the conditional execution of so many succeeding instructions using the execution control instruction.

SUMMARY OF THE INVENTION

An object of the present invention is controlling the conditional execution of as many succeeding instructions as possible using an execution control instruction of a short word length to suppress the branch hazard in an information processor for processing the instructions by pipelining.

To achieve this object, the present invention introduces an instruction-specifying field, which is used to define, in binary code, the number of instructions that should be executed under controlled conditions, into the execution control instruction.

More specifically, the present invention uses an execution control instruction, which contains: a condition field specifying an execution condition; and an instruction-specifying field defining, in binary code, the number of instructions to be executed only conditionally. A decision is made as to whether or not the execution condition that has been specified by the condition field is satisfied. Based on the outcome of this decision, it is determined whether or not that number of instructions, which has been defined by the instruction-specifying field for instructions succeeding the execution control instruction, should be nullified.

BRIEF DESCRIPTION OF THE DRAWINGS

Figure 1 is a block diagram schematically illustrating

an exemplary configuration for a program execution control system according to the present invention.

Figure 2 illustrates a basic format for an execution control instruction according to the present invention.

5 Figure 3 illustrates a format for an execution control instruction according to a first embodiment of the present invention.

Figure 4 is a flowchart illustrating a program execution control procedure according to the first embodiment.

10 Figures 5A and 5B illustrate specific examples of the program execution control according to the first embodiment.

Figure 6 illustrates a format for an execution control instruction according to a second embodiment of the present invention.

15 Figure 7 is a flowchart illustrating a program execution control procedure according to the second embodiment.

Figures 8A and 8B illustrate specific examples of the program execution control according to the second embodiment.

20 Figure 9 illustrates a format for an execution control instruction according to a third embodiment of the present invention.

Figure 10 is a flowchart illustrating a program execution control procedure according to the third embodiment.

25 Figures 11A and 11B illustrate specific examples of the

program execution control according to the third embodiment.

Figure 12 illustrates a format for an execution control instruction according to a fourth embodiment of the present invention.

5 Figure 13 is a flowchart illustrating a program execution control procedure according to the fourth embodiment.

10 Figures 14A, 14B, 14C and 14D illustrate specific examples of the program execution control according to the fourth embodiment.

Figure 15 illustrates a format for an execution control instruction according to a fifth embodiment of the present invention.

15 Figure 16 is a flowchart illustrating a program execution control procedure according to the fifth embodiment.

Figures 17A, 17B, 17C and 17D illustrate specific examples of the program execution control according to the fifth embodiment.

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DETAILED DESCRIPTION OF THE INVENTION

Figure 1 schematically illustrates an exemplary configuration for a program execution control system according to the present invention. The system shown in
25 Figure 1 is used for controlling the execution of a program

in an information processor of the type processing instructions by pipelining. As shown in Figure 1, the system includes instruction providing section 10, instruction executing section 20 and nullification controller 30. The instruction providing section 10 includes program counter 11, memory 12 and instruction register 13. The program counter 11 sequentially specifies respective addresses of instructions to be fetched and executed. The memory 12 reads out the instructions, which together constitute a program including an execution control instruction, one after another in accordance with the addresses specified. And the instruction register 13 stores thereon the instructions read out one by one. An instruction set **INST** is provided from the instruction register 13 to the instruction executing section 20 and to the nullification controller 30. The instruction executing section 20 includes: an arithmetic logic unit (ALU) 21 for performing arithmetic or logic operations in response to the instruction set **INST** provided; and a flag register 22 for holding multiple flags indicating the results of the computations. The nullification controller 30 receives control flags **CF** from the flag register 22. In the illustrated embodiment, the control flags **CF** include two flags **F1** and **F2**, which have been set equal to zero or one by the ALU 21.

Figure 2 illustrates a basic format for an execution control instruction according to the present invention. As

shown in Figure 2, the execution control instruction contains instruction field **41**, condition field **42** and instruction-specifying field **43**. The instruction field **41** specifies the type of the instruction as an execution control instruction.

5 The condition field **42** specifies an execution condition **EC**. And the instruction-specifying field **43** defines, in binary code, the number **N** (where **N** is a natural number) of instructions that will be executed only conditionally, i.e., just when the execution condition **EC** is satisfied. In

10 response to the execution control instruction, the nullification controller **30** decides based on the control flags **CF** whether or not the execution condition **EC** specified by the condition field **42** is met. And based on the outcome of this decision, the nullification controller **30** determines

15 whether or not the number **N** of instructions, which number has been defined by the instruction-specifying field **43** for instructions succeeding the execution control instruction, should be nullified. Suppose the nullification controller **30** has determined that the number **N** of instructions should be

20 nullified since the execution condition **EC** is not met, the controller **30** asserts a nullification signal **NUL** to be supplied to the ALU **21**. In that case, even if the number **N** of instructions following the execution control instruction have already been executed to a midway point of the pipeline,

25 these instructions are nullified in the ALU **21**. Accordingly,

the same results are attained as if NOP (no operation) instructions had been executed instead of these instructions. Alternatively, if the nullification controller 30 has determined that the number N of succeeding instructions should not be nullified since the execution condition EC is met, then the nullification signal NUL is not asserted. Thus, the number N of instructions succeeding the execution control instruction are enabled and executed by the ALU 21.

Hereinafter, preferred embodiments of the present invention will be described with reference to Figures 3 through 17.

EMBODIMENT 1

Figure 3 illustrates a 16-bit format of an execution control instruction according to a first embodiment of the present invention. As shown in Figure 3, the execution control instruction contains a 12-bit instruction field 41, a 2-bit condition field 42 and a 2-bit instruction-specifying field 43. The instruction field 41 sets an operation code OP specifying the instruction type of the operation to be performed, i.e., execution control instruction. The condition field 42 specifies any of four types of execution conditions. In the illustrated embodiment, the execution condition EC to be specified is a binary code representing "F1 = 1", "F1 = 0", "F2 = 1" or "F2 = 0". The instruction-

specifying field 43 specifies, in binary code, the number N (where N=1 through 4) of instructions that will be executed under controlled conditions. In the illustrated embodiment, codes "00", "01", "10" and "11" represent "N=1", "N=2", "N=3" and "N=4", respectively.

Figure 4 illustrates a program execution control procedure according to the first embodiment. First, in Step S1, the execution condition EC and the number N of instructions to be executed under controlled conditions are defined for the condition field 42 and instruction-specifying field 43, respectively. Next, in Step S2, the execution control instruction is executed. Then, in Step S3, it is determined based on the control flags CF whether or not the execution condition EC is satisfied. If the execution condition EC is not met, then the number N of instructions succeeding the execution control instruction are nullified in Step S4. Alternatively, if the execution condition EC is met, then the number N of instructions following the execution control instruction are executed in Step S5.

Figures 5A and 5B illustrate specific examples of the program execution control according to the first embodiment. In the illustrated examples, "F1 = 1" is the execution condition EC and "N = 4". As shown in Figure 5A, if the execution condition is not met (i.e., when F1=0), then the four instructions succeeding the execution control

instruction (i.e., Succeeding Instructions 1 through 4) are nullified. On the other hand, if the execution condition is met (i.e., when $F1=1$), then the four instructions following the execution control instruction (i.e., Succeeding Instructions 1 through 4) are executed, not nullified, as shown in Figure 5B.

In the first embodiment, the number of bits in the condition field 42 is two. Optionally, the number of execution conditions may be increased by increasing this bit number. The number of bits in the instruction-specifying field 43 is also two in the first embodiment. Alternatively, the number of instructions to be executed under controlled conditions may be increased by increasing this bit number. These modified examples will be equally applicable to all of the following second through fifth embodiments.

EMBODIMENT 2

Figure 6 illustrates a 16-bit format for an execution control instruction according to a second embodiment of the present invention. The format shown in Figure 6 is the same as that shown in Figure 3.

Figure 7 illustrates a program execution control procedure according to the second embodiment. First, in Step S11, the execution condition EC and the number N of instructions to be executed under controlled conditions are

defined for the condition field 42 and instruction-specifying field 43, respectively. Next, in Step S12, the execution control instruction is executed. Then, in Step S13, it is determined based on the control flags CF whether or not the execution condition EC is satisfied. If the execution condition EC is not met, then the number N of instructions succeeding the execution control instruction are nullified in Step S14a. Thereafter, the next number N of instructions following the former instruction set are executed in Step S14b. Alternatively, if the execution condition EC is met, then the number N of instructions succeeding the execution control instruction are executed in Step S15a. Thereafter, the next number N of instructions following the former instruction set are nullified in Step S15b.

Figures 8A and 8B illustrate specific examples of the program execution control according to the second embodiment. In the illustrated examples, "F1 = 1" is the execution condition EC and "N = 2". As shown in Figure 8A, if the execution condition is not met (i.e., when F1 = 0), then the first two instructions succeeding the execution control instruction (i.e., Succeeding Instructions 1 and 2) are nullified, and the next two instructions (i.e., Succeeding Instructions 3 and 4) are executed. On the other hand, if the execution condition is met (i.e., when F1 = 1), then the first two instructions following the execution control

instruction (i.e., Succeeding Instructions 1 and 2) are executed, and the next two instructions (i.e., Succeeding Instructions 3 and 4) are nullified as shown in Figure 8B. That is to say, according to this embodiment, instructions can be executed conditionally in accordance with the if-then-else construct.

EMBODIMENT 3

Figure 9 illustrates a 16-bit format for an execution control instruction according to a third embodiment of the present invention. As shown in Figure 9, the execution control instruction contains a 10-bit instruction field 41, a 2-bit condition field 42 and a 4-bit instruction-specifying field 43. The instruction field 41 defines an operation code OP specifying the instruction type of the operation to be performed, i.e., execution control instruction. The condition field 42 specifies any of four types of execution conditions. The instruction-specifying field 43 consists of first and second instruction-specifying sub-fields 43a and 43b, each defining, in binary code, the number of instructions that will be executed only conditionally. The first instruction-specifying sub-field 43a is a 2-bit field for defining, in binary code, the first number N1 (where N1=1 through 4) of instructions to be executed under controlled conditions. The second instruction-specifying sub-field 43b is also a 2-bit

field for defining, in binary code, the second number **N2** (where **N2**=1 through 4) of instructions to be executed under controlled conditions.

Figure 10 illustrates a program execution control procedure according to the third embodiment. First, in Step **S21**, the execution condition **EC** and the first and second numbers **N1** and **N2** of instructions to be executed under controlled conditions are defined for the condition field **42** and instruction-specifying sub-fields **43a** and **43b**, respectively. Next, in Step **S22**, the execution control instruction is executed. Then, in Step **S23**, it is determined based on the control flags **CF** whether or not the execution condition **EC** is satisfied. If the execution condition **EC** is not met, then the first number **N1** of instructions succeeding the execution control instruction are nullified in Step **S24a**. Thereafter, the second number **N2** of instructions following the first instruction set are executed in Step **S24b**. Alternatively, if the execution condition **EC** is met, then the first number **N1** of instructions succeeding the execution control instruction are executed in Step **S25a**. Thereafter, the second number **N2** of instructions following the first instruction set are nullified in Step **S25b**.

Figures 11A and 11B illustrate specific examples of the program execution control according to the third embodiment. In the illustrated examples, "**F1** = 1" is the execution

condition **EC**, "N1=1" and "N2=3". As shown in Figure 11A, if the execution condition is not met (i.e., when F1=0), then the first instruction succeeding the execution control instruction (i.e., Succeeding Instruction 1) is nullified, and the next three instructions (i.e., Succeeding Instructions 2, 3 and 4) are executed. On the other hand, if the execution condition is met (i.e., when F1=1), then the first instruction following the execution control instruction (i.e., Succeeding Instruction 1) is executed, and the next three instructions (i.e., Succeeding Instructions 2, 3 and 4) are nullified as shown in Figure 11B. That is to say, according to this embodiment, instructions can be executed conditionally in accordance with the if-then-else construct. In addition, the respective numbers of instructions corresponding to the THEN and ELSE statements can be defined independently.

EMBODIMENT 4

Figure 12 illustrates a 16-bit format for an execution control instruction according to a fourth embodiment of the present invention. As shown in Figure 12, the execution control instruction contains a 10-bit instruction field **41**, a 4-bit condition field **42** and a 2-bit instruction-specifying field **43**. The instruction field **41** defines an operation code **OP** specifying the instruction type of the operation to be

performed, i.e., execution control instruction. The condition field 42 consists of first and second condition sub-fields 42a and 42b, each specifying a single execution condition. The first condition sub-field 42a is a 2-bit field for specifying any of four types of execution conditions as a first execution condition EC1. The second condition sub-field 42b is also a 2-bit field for specifying any of four types of execution conditions as a second execution condition EC2. The instruction-specifying sub-field 43 defines, in binary code, the number N (where N = 1 through 4) of instructions to be executed under controlled conditions.

Figure 13 illustrates a program execution control procedure according to the fourth embodiment. First, in Step S31, the first and second execution conditions EC1 and EC2 and the number N of instructions to be executed under controlled conditions are defined for the first and second condition sub-fields 42a and 42b and instruction-specifying field 43, respectively. Next, in Step S32, the execution control instruction is executed. Then, in Step S33a, it is determined based on the control flags CF whether or not the first execution condition EC1 is satisfied. Next, in Step S33b, it is determined based on the control flags CF whether or not the second execution condition EC2 is satisfied. If neither the first nor second execution condition EC1 nor EC2 is met, then the number N of instructions succeeding the execution control

instruction are nullified in Step **S34a**. Thereafter, the next
number **N** of instructions succeeding the former instruction set
are nullified in Step **S34b**. Alternatively, if the first
execution condition **EC1** is not met and the second execution
5 condition **EC2** is met, then the number **N** of instructions
succeeding the execution control instruction are nullified in
Step **S35a**. Thereafter, the next number **N** of instructions
succeeding the former instruction set are executed in Step
S35b. As another alternative, if the first execution condition
10 **EC1** is met and the second execution condition **EC2** is not met,
then the number **N** of instructions succeeding the execution
control instruction are executed in Step **S36a**. Thereafter,
the next number **N** of instructions succeeding the former
instruction set are nullified in Step **S36b**. As further
15 alternative, if both the first and second execution conditions
EC1 and **EC2** are met, then the number **N** of instructions
succeeding the execution control instruction are executed in
Step **S37a**. Thereafter, the next number **N** of instructions
succeeding the former instruction set are executed in Step
20 **S37b**.

Figures **14A** through **14D** illustrate specific examples of
the program execution control according to the fourth
embodiment. In the illustrated examples, "**F1=1**" is the first
execution condition **EC1**, "**F2 = 0**" is the second execution
25 condition **EC2** and "**N=2**". As shown in Figure **14A**, if neither

the first nor second execution condition is met (i.e., when $F1 = 0$ and $F2 = 1$), then four instructions succeeding the execution control instruction (i.e., Succeeding Instructions 1 through 4) are nullified. Alternatively, if the first execution condition is not met and the second execution condition is met (i.e., when $F1 = 0$ and $F2 = 0$), then two instructions succeeding the execution control instruction (i.e., Succeeding Instructions 1 and 2) are nullified and next two instructions succeeding the former instruction set (i.e., Succeeding Instructions 3 and 4) are executed as shown in Figure 14B. As another alternative, if the first execution condition is met and the second execution condition is not met (i.e., when $F1 = 1$ and $F2 = 1$), then two instructions succeeding the execution control instruction (i.e., Succeeding Instructions 1 and 2) are executed and next two instructions succeeding the former instruction set (i.e., Succeeding Instructions 3 and 4) are nullified as shown in Figure 14C. As further alternative, if both the first and second execution conditions are met (i.e., when $F1 = 1$ and $F2 = 0$), then four instructions succeeding the execution control instruction (i.e., Succeeding Instructions 1 through 4) are executed, not nullified, as shown in Figure 14D.

It should be noted that the number M of condition sub-fields in the execution control instruction may be equal to or larger than three. In such a case, a number N of

instructions, which number has been defined by the instruction-specifying field **43** for instructions succeeding the execution control instruction, are regarded as instructions to be executed under controlled conditions. If
5 an execution condition **EC_m** (where $m=1$ through M), which has been specified by associated one of the number M of condition sub-fields, is not met, then the step of nullifying the number N of instructions at a location corresponding to the specified execution condition **EC_m** is performed. Furthermore,
10 if an execution condition **EC_m** (where $m=1$ through M), which has been specified by associated one of the number M of condition sub-fields, is met, then the step of executing the number N of instructions at a location corresponding to the specified execution condition **EC_m** is performed.

EMBODIMENT 5

Figure 15 illustrates a 16-bit format for an execution control instruction according to a fifth embodiment of the present invention. As shown in Figure 15, the execution
20 control instruction contains an 8-bit instruction field **41**, a 4-bit condition field **42** and a 4-bit instruction-specifying field **43**. The instruction field **41** defines an operation code **OP** specifying the instruction type of the operation to be performed, i.e., execution control instruction. The condition
25 field **42** consists of first and second condition sub-fields

42a and 42b, each specifying a single execution condition. The first condition sub-field 42a is a 2-bit field for specifying any of four types of execution conditions as a first execution condition EC1. The second condition sub-field 42b is also a 2-bit field for specifying any of four types of execution conditions as a second execution condition EC2. The instruction-specifying field 43 consists of first and second instruction-specifying sub-fields 43a and 43b, each defining, in binary code, the number of conditionally executable instructions. The first instruction-specifying sub-field 43a is a 2-bit field for defining, in binary code, the first number N1 (where N1=1 through 4) of instructions to be executed under controlled conditions. The second instruction-specifying sub-field 43b is also a 2-bit field for defining, in binary code, the second number N2 (where N2 = 1 through 4) of instructions to be executed under controlled conditions.

Figure 16 illustrates a program execution control procedure according to the fifth embodiment. First, in Step S41, the first and second execution conditions EC1 and EC2 and first and second numbers N1 and N2 of conditionally executable instructions are defined for the first and second condition sub-fields 42a and 42b and first and second instruction-specifying sub-fields 43a and 43b, respectively. Next, in Step S42, the execution control instruction is executed. Then, in

Step **S43a**, it is determined based on the control flags **CF** whether or not the first execution condition **EC1** is satisfied. Next, in Step **S43b**, it is determined based on the control flags **CF** whether or not the second execution condition **EC2** is satisfied. If neither the first nor second execution condition **EC1** nor **EC2** is met, then the number **N1** of instructions succeeding the execution control instruction are nullified in Step **S44a**. Thereafter, the next number **N2** of instructions succeeding the former instruction set are nullified in Step **S44b**. Alternatively, if the first execution condition **EC1** is not met and the second execution condition **EC2** is met, then the number **N1** of instructions succeeding the execution control instruction are nullified in Step **S45a**. Thereafter, the next number **N2** of instructions succeeding the former instruction set are executed in Step **S45b**. As another alternative, if the first execution condition **EC1** is met and the second execution condition **EC2** is not met, then the number **N1** of instructions succeeding the execution control instruction are executed in Step **S46a**. Thereafter, the next number **N2** of instructions succeeding the former instruction set are nullified in Step **S46b**. As further alternative, if both the first and second execution conditions **EC1** and **EC2** are met, then the number **N1** of instructions succeeding the execution control instruction are executed in Step **S47a**. Thereafter, the next number **N2** of instructions succeeding the

former instruction set are executed in Step **S47b**.

Figures **17A** through **17D** illustrate specific examples of the program execution control according to the fifth embodiment. In the illustrated examples, "**F1 = 1**" is the first execution condition **EC1**, "**F2 = 0**" is the second execution condition **EC2**, "**N1 = 1**" and "**N2 = 3**". As shown in Figure **17A**, if neither the first nor second execution condition is met (i.e., when **F1 = 0** and **F2 = 1**), then four instructions succeeding the execution control instruction (i.e., Succeeding Instructions 1 through 4) are nullified. Alternatively, if the first execution condition is not met and the second execution condition is met (i.e., when **F1 = 0** and **F2 = 0**), then one instruction succeeding the execution control instruction (i.e., Succeeding Instruction 1) is nullified and the next three instructions following the former instruction (i.e., Succeeding Instructions 2, 3 and 4) are executed as shown in Figure **17B**. As another alternative, if the first execution condition is met and the second execution condition is not met (i.e., when **F1 = 1** and **F2 = 1**), then one instruction succeeding the execution control instruction (i.e., Succeeding Instruction 1) is executed and the next three instructions succeeding the former instruction (i.e., Succeeding Instructions 2, 3 and 4) are nullified as shown in Figure **17C**. As further alternative, if both the first and second execution conditions are met (i.e., when **F1 = 1** and **F2 =**

0), then four instructions succeeding the execution control instruction (i.e., Succeeding Instructions 1 through 4) are executed, not nullified, as shown in Figure 17D.

It should be noted that the number **M** of condition sub-
5 fields of the execution control instruction may be equal to or larger than three. In such a case, the same number **M** of instruction-specifying sub-fields, which correspond to the number **M** of condition sub-fields, respectively, are included in the instruction-specifying field of the execution control
10 instruction. A number **Nm** of instructions (where $m=1$ through **M**), which number has been defined by associated one of the number **M** of instruction-specifying sub-fields for instructions succeeding the execution control instruction, are regarded as instructions to be executed under controlled
15 conditions. If an execution condition **ECm** (where $m = 1$ through **M**), which is specified by associated one of the number **M** of condition sub-fields, is not met, then the step of nullifying the number **Nm** of instructions at a location corresponding to the specified execution condition **ECm** is
20 performed. Furthermore, if the execution condition **ECm** (where $m = 1$ through **M**), which has been specified by associated one of the number **M** of condition sub-fields, is met, then the step of executing the number **Nm** of instructions at a location corresponding to the specified execution
25 condition **ECm** is performed.

WHAT IS CLAIMED IS:

1. A system for controlling the execution of a program in an information processor for processing instructions by pipelining, the system comprising:

means for providing a set of instructions including an execution control instruction, the execution control instruction containing a condition field and an instruction-specifying field, the condition field specifying an execution condition, the instruction-specifying field defining, in binary code, the number of instructions to be executed only conditionally;

means for deciding whether or not the execution condition that has been specified by the condition field is satisfied; and

means for determining based on the outcome of the decision whether or not said number of instructions, which number has been defined by the instruction-specifying field for instructions succeeding the execution control instruction, should be nullified.

2. A method for controlling the execution of a program in an information processor for processing instructions by pipelining, the method comprising the steps of:

a) providing an execution control instruction, the execution control instruction containing a condition field

and an instruction-specifying field, the condition field specifying an execution condition, the instruction-specifying field defining, in binary code, the number of instructions to be executed only conditionally;

b) deciding whether or not the execution condition that has been specified by the condition field is satisfied; and

c) determining based on the outcome of the decision step b) whether or not said number of instructions, which number has been defined by the instruction-specifying field for instructions succeeding the execution control instruction, should be nullified.

3. The method of claim 2, wherein the condition field is a single field for specifying the execution condition, and

wherein the instruction-specifying field is a single field for defining the instruction number, and

wherein the step c) comprises the sub-step of regarding said number of instructions, which number has been defined by the instruction-specifying field for instructions succeeding the execution control instruction, as instructions to be executed only conditionally, and nullifying the conditionally executable instructions if the execution condition that has been specified by the condition field is not satisfied.

4. The method of claim 3, wherein the step c) further

comprises the sub-step of executing the conditionally executable instructions if the execution condition that has been specified by the condition field is satisfied.

5. The method of claim 2, wherein the condition field is a single field for specifying the execution condition, and

wherein the instruction-specifying field is a single field for defining the instruction number, and

wherein the step c) comprises the sub-steps of:

regarding said number of instructions, which number has been defined by the instruction-specifying field for instructions succeeding the execution control instruction, as a first set of conditionally executable instructions, and nullifying the first set of conditionally executable instructions if the execution condition that has been specified by the condition field is not satisfied; and

regarding said number of instructions, which number has been defined by the instruction-specifying field for instructions succeeding the first set of conditionally executable instructions, as a second set of conditionally executable instructions, and nullifying the second set of conditionally executable instructions if the execution condition that has been specified by the condition field is satisfied.

6. The method of claim 5, wherein the step c) further comprises the sub-steps of:

getting the first set of conditionally executable instructions executed if the execution condition that has been specified by the condition field is satisfied; and

getting the second set of conditionally executable instructions executed if the execution condition that has been specified by the condition field is not satisfied.

7. The method of claim 2, wherein the condition field is a single field for specifying the execution condition, and

wherein the instruction-specifying field contains first and second instruction-specifying sub-fields, which respectively define, in binary code, first and second numbers of instructions to be executed only conditionally, and

wherein the step c) comprises the sub-steps of:

regarding the first number of instructions, which number has been defined by the first instruction-specifying sub-field for instructions succeeding the execution control instruction, as a first set of conditionally executable instructions, and nullifying the first set of conditionally executable instructions if the execution condition specified by the condition field is not satisfied; and

regarding the second number of instructions, which number has been defined by the second instruction-specifying

sub-field for instructions succeeding the first set of conditionally executable instructions, as a second set of conditionally executable instructions, and nullifying the second set of conditionally executable instructions if the execution condition specified by the condition field is satisfied.

8. The method of claim 7, wherein the step c) further comprises the sub-steps of:

getting the first set of conditionally executable instructions executed if the execution condition that has been specified by the execution field is satisfied; and

getting the second set of conditionally executable instructions executed if the execution condition that has been specified by the execution field is not satisfied.

9. The method of claim 2, wherein the condition field includes multiple condition sub-fields, each specifying a single associated execution condition, and

wherein the instruction-specifying field is a single field for defining the instruction number, and

wherein the step c) comprises a plurality of sub-steps, and

wherein in each said sub-step, said number of instructions, which number has been defined by the

instruction-specifying field for instructions succeeding the execution control instruction, are regarded as conditionally executable instructions, and, if the execution condition specified by associated one of the condition sub-fields is not satisfied, the conditionally executable instructions at a location corresponding to the execution condition specified are nullified.

10. The method of claim 9, wherein the step c) comprises a plurality of sub-steps,

wherein in each said sub-step, if the execution condition specified by associated one of the condition sub-fields is satisfied, the conditionally executable instructions at a location corresponding to the execution condition specified are executed.

11. The method of claim 2, wherein the condition field includes multiple condition sub-fields, each specifying a single execution condition, and

wherein the instruction-specifying field includes multiple instruction-specifying sub-fields corresponding to the respective condition sub-fields, each said instruction-specifying sub-field defining an associated instruction number in binary code, and

wherein the step c) comprises a plurality of sub-steps,

and

wherein in each said sub-step, said number of instructions, which number has been defined by associated one of the instruction-specifying sub-fields for instructions succeeding the execution control instruction, are regarded as conditionally executable instructions, and, if the execution condition specified by associated one of the condition sub-fields is not satisfied, the conditionally executable instructions at a location corresponding to the execution condition specified are nullified.

12. The method of claim 11, wherein the step c) comprises a plurality of sub-steps,

wherein in each said sub-step, if the execution condition specified by associated one of the condition sub-fields is satisfied, the conditionally executable instructions at the location corresponding to the execution condition specified are executed.

ABSTRACT OF THE DISCLOSURE

An execution control instruction is applied to an information processor of the type processing instructions by pipelining to suppress the occurrence of branch hazard. The execution control instruction contains: a condition field for specifying an execution condition; and an instruction-specifying field for defining, in binary code, the number of instructions to be executed only conditionally. In response to the execution control instruction, a nullification controller decides, based on control flags provided from an arithmetic logic unit, whether or not the execution condition specified by the condition field is satisfied. And based on the outcome of this decision, the controller determines whether or not that number of instructions, which has been defined by the instruction-specifying field for instructions succeeding the execution control instruction, should be nullified. If the controller has determined that the specified number of succeeding instructions should be nullified since the execution condition is not met, then the controller asserts a nullification signal to be supplied to the arithmetic logic unit. In this manner, a large number of succeeding instructions are executable conditionally using an execution control instruction of a short word length.

Fig. 1

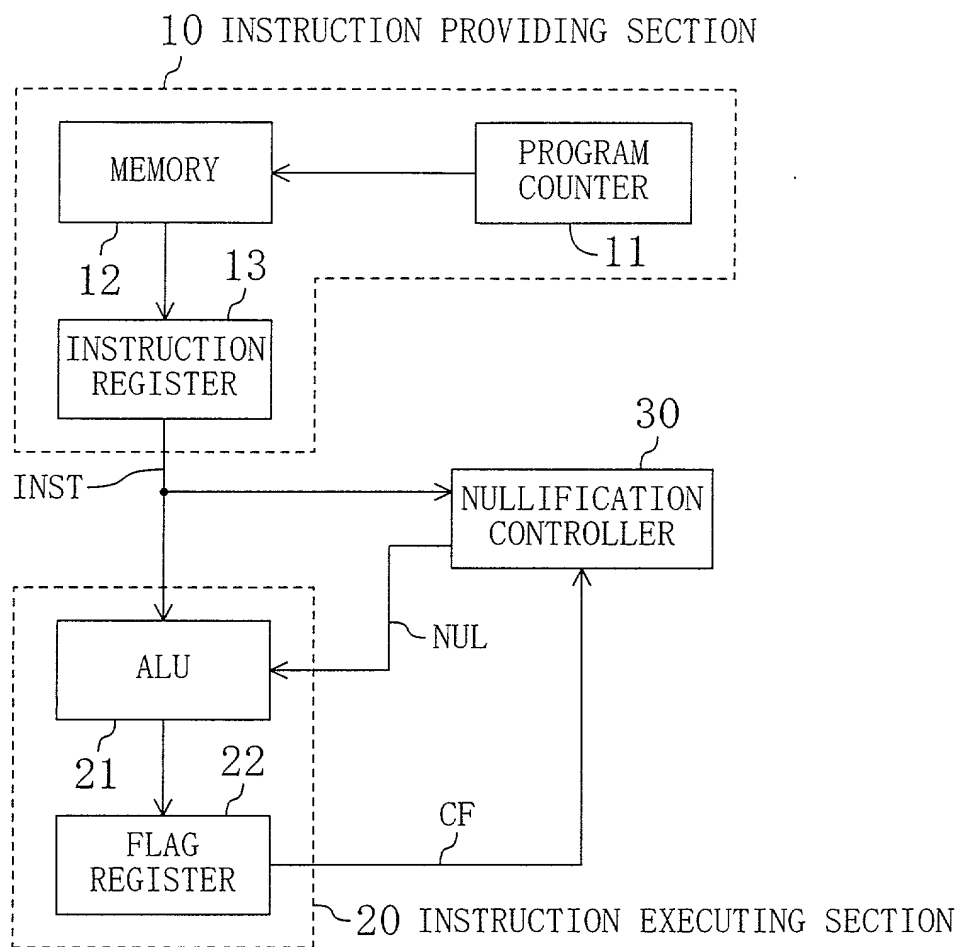


Fig. 2

EXECUTION CONTROL INSTRUCTION

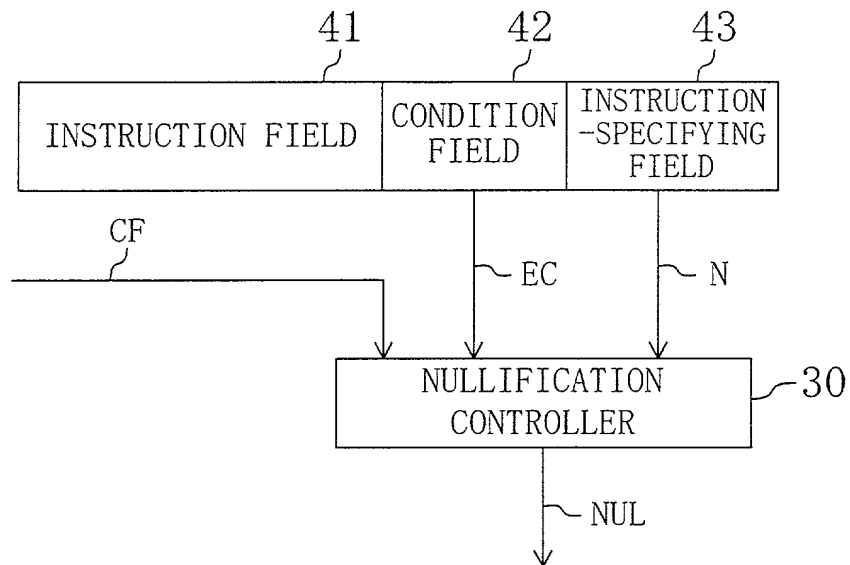


Fig. 3

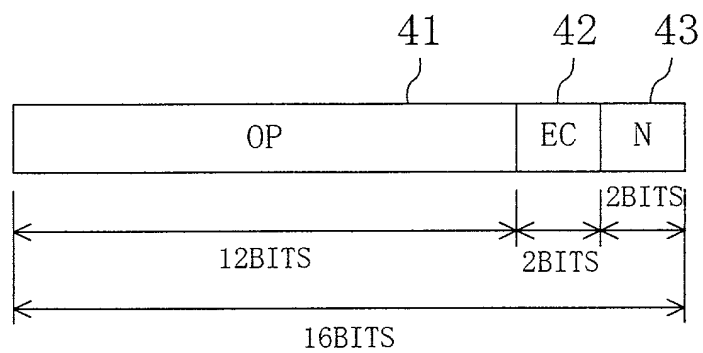


Fig. 4

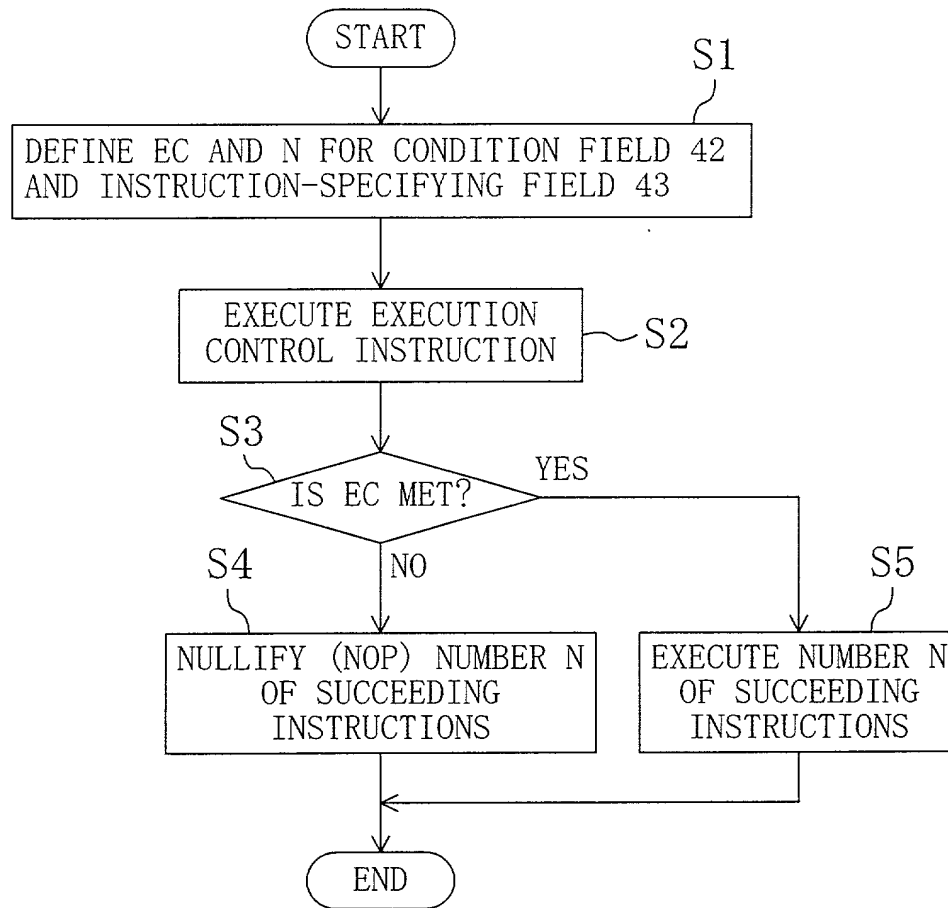


Fig. 15

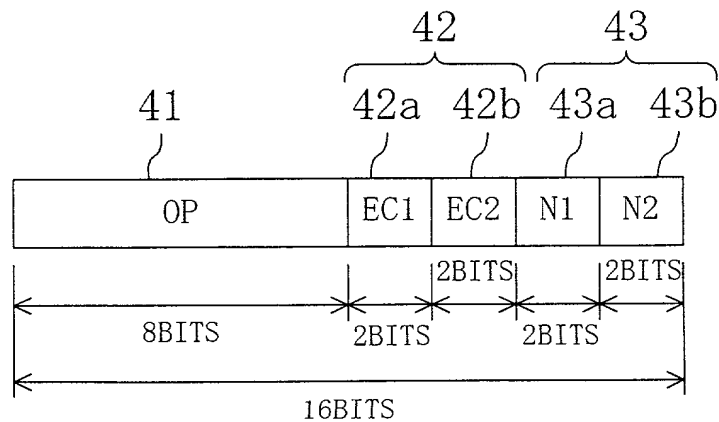


Fig. 5A

IF EC IS NOT MET ($F1=0$)

EXECUTION CONTROL INSTRUCTION F1, 4

NOP

NOP

NOP

NOP

SUCCEEDING INSTRUCTION 5

Fig. 5B

IF EC IS MET ($F1=1$)

EXECUTION CONTROL INSTRUCTION F1, 4

SUCCEEDING INSTRUCTION 1

SUCCEEDING INSTRUCTION 2

SUCCEEDING INSTRUCTION 3

SUCCEEDING INSTRUCTION 4

SUCCEEDING INSTRUCTION 5

Fig. 6

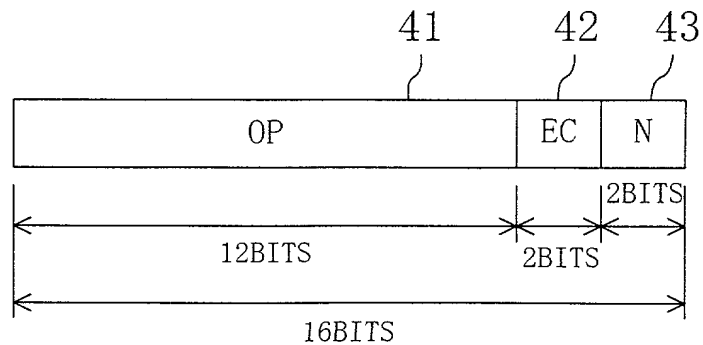


Fig. 7

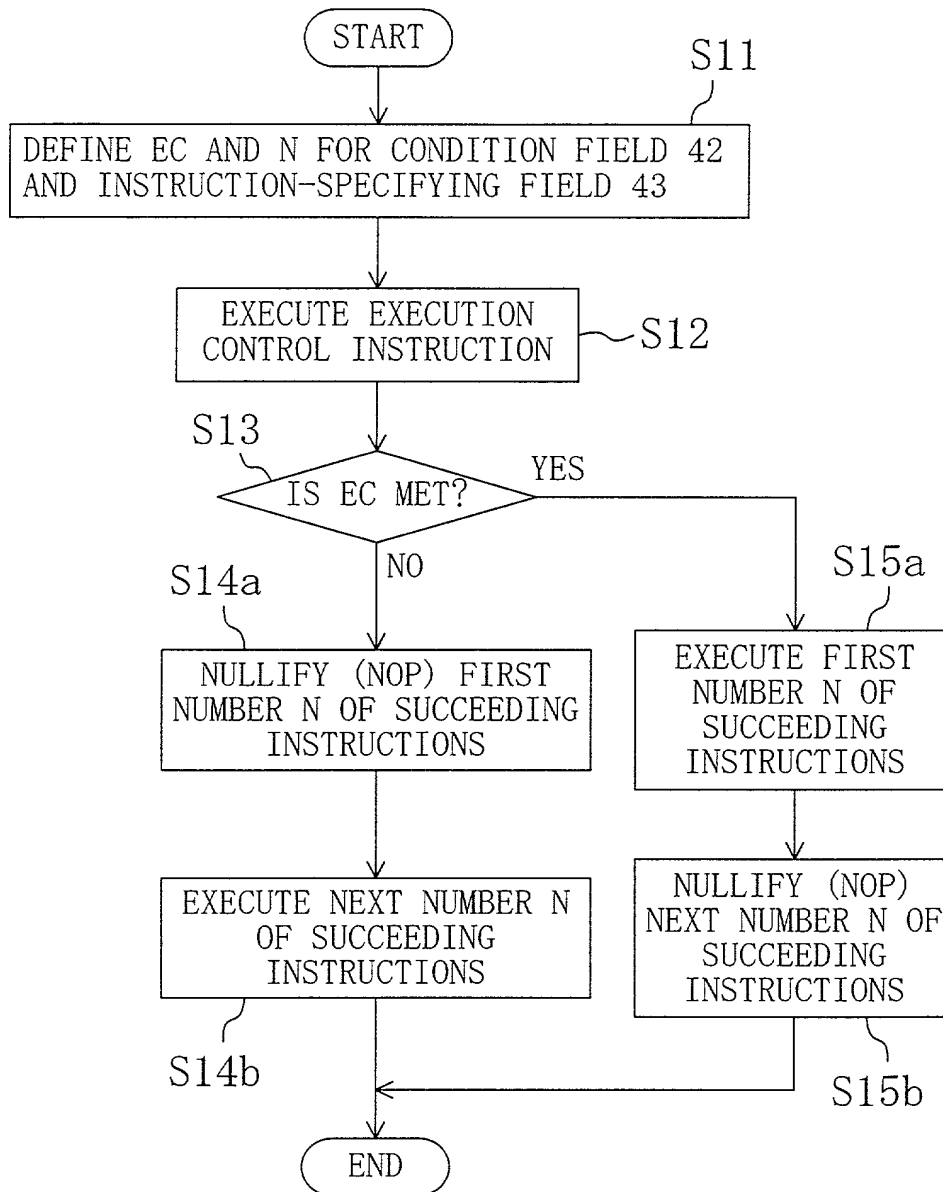


Fig. 8A

IF EC IS NOT MET ($F1=0$)

EXECUTION CONTROL INSTRUCTION F1, 2

NOP

NOP

SUCCEEDING INSTRUCTION 3

SUCCEEDING INSTRUCTION 4

SUCCEEDING INSTRUCTION 5

Fig. 8B

IF EC IS MET ($F1=1$)

EXECUTION CONTROL INSTRUCTION F1, 2

SUCCEEDING INSTRUCTION 1

SUCCEEDING INSTRUCTION 2

NOP

NOP

SUCCEEDING INSTRUCTION 5

Fig. 9

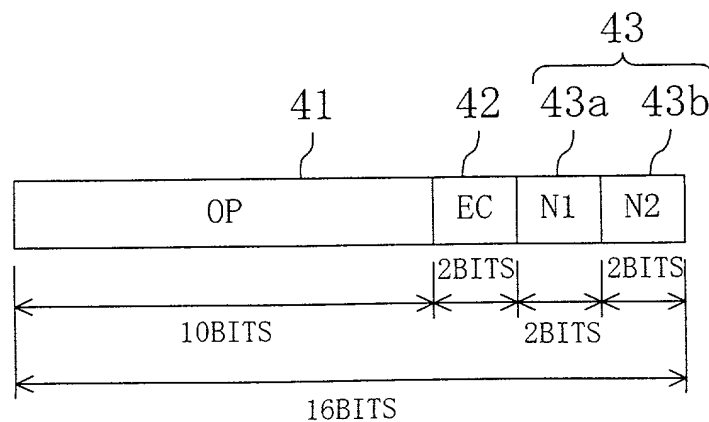


Fig. 10

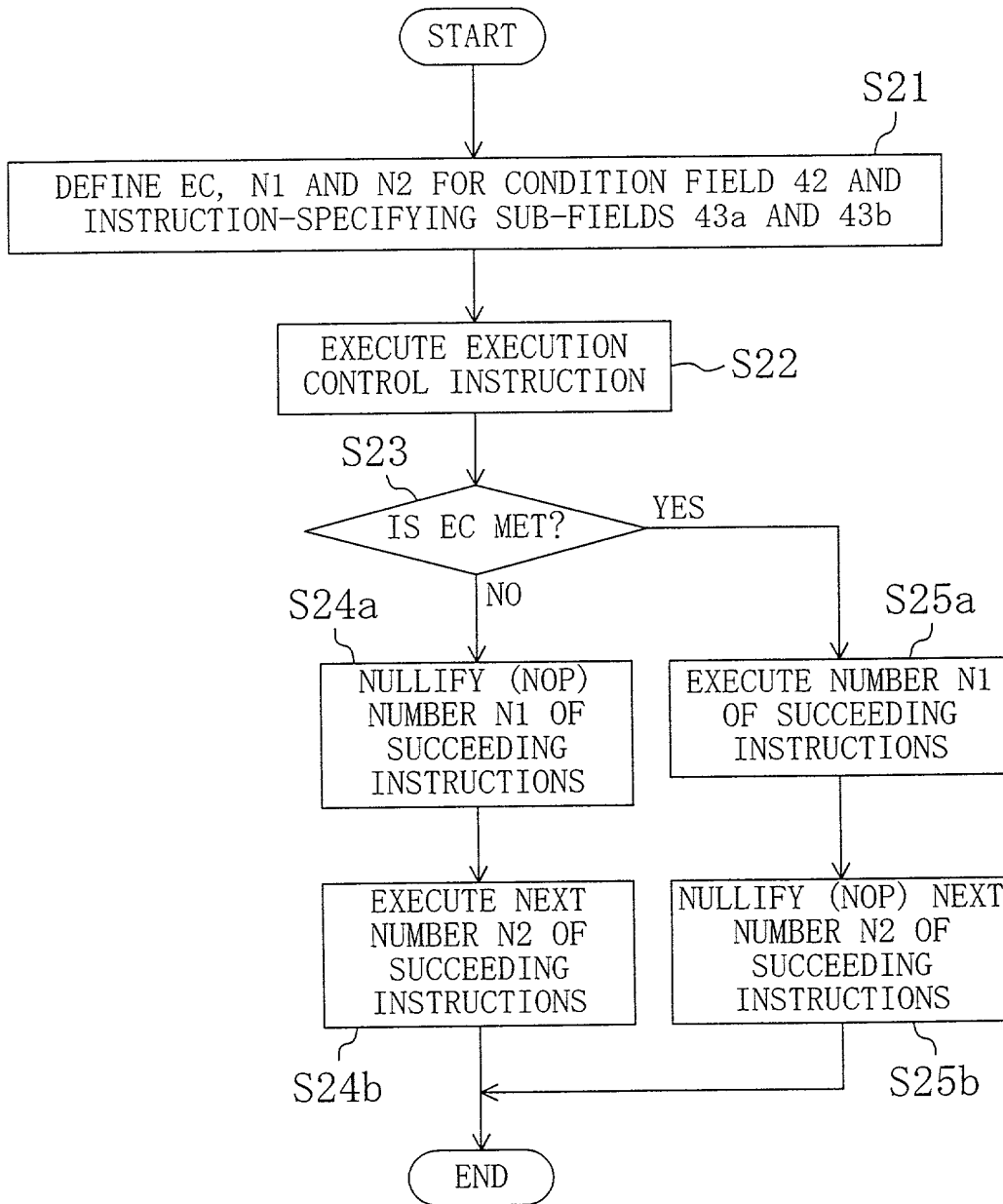


Fig. 11A

IF EC IS NOT MET ($F1=0$)

EXECUTION CONTROL INSTRUCTION F1, 1, 3
NOP

SUCCEEDING INSTRUCTION 2

SUCCEEDING INSTRUCTION 3

SUCCEEDING INSTRUCTION 4

SUCCEEDING INSTRUCTION 5

Fig. 11B

IF EC IS MET ($F1=1$)

EXECUTION CONTROL INSTRUCTION F1, 1, 3
SUCCEEDING INSTRUCTION 1

NOP

NOP

NOP

SUCCEEDING INSTRUCTION 5

Fig. 12

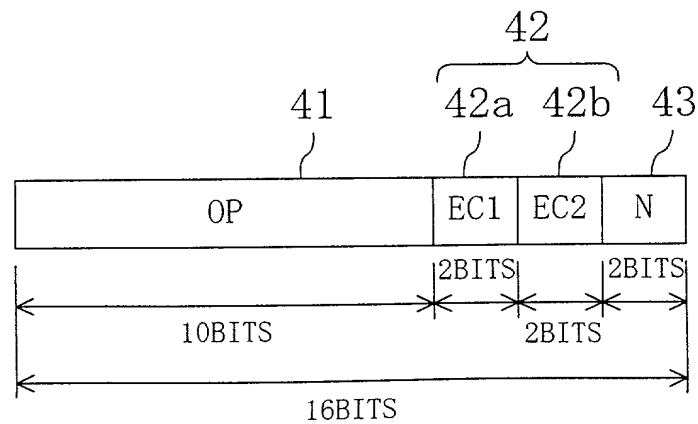


Fig. 13

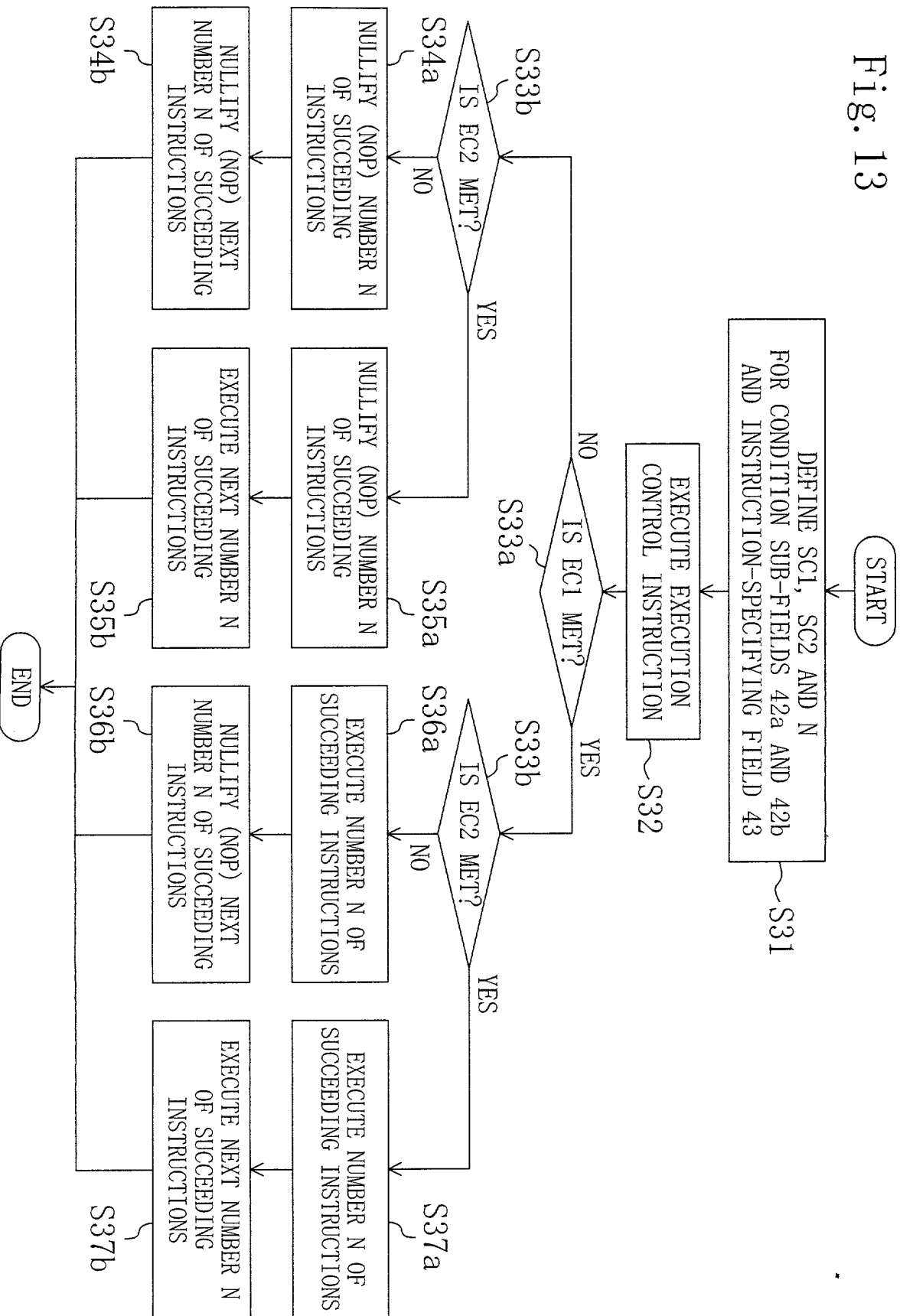


Fig. 14A

IF NEITHER EC1 NOR EC2 IS MET (WHEN $F1=0$ AND $F2=1$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 2$
NOP
NOP
NOP
NOP
SUCCEEDING INSTRUCTION 5

Fig. 14B

IF EC1 IS NOT MET AND EC2 IS MET (WHEN $F1=0$ AND $F2=0$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 2$
NOP
NOP
SUCCEEDING INSTRUCTION 3
SUCCEEDING INSTRUCTION 4
SUCCEEDING INSTRUCTION 5

Fig. 14C

IF EC1 IS MET AND EC2 IS NOT MET (WHEN $F1=1$ AND $F2=1$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 2$
SUCCEEDING INSTRUCTION 1
SUCCEEDING INSTRUCTION 2
NOP
NOP
SUCCEEDING INSTRUCTION 5

Fig. 14D

IF BOTH EC1 AND EC2 ARE MET (WHEN $F1=1$ AND $F2=0$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 2$
SUCCEEDING INSTRUCTION 1
SUCCEEDING INSTRUCTION 2
SUCCEEDING INSTRUCTION 3
SUCCEEDING INSTRUCTION 4
SUCCEEDING INSTRUCTION 5

Fig. 16

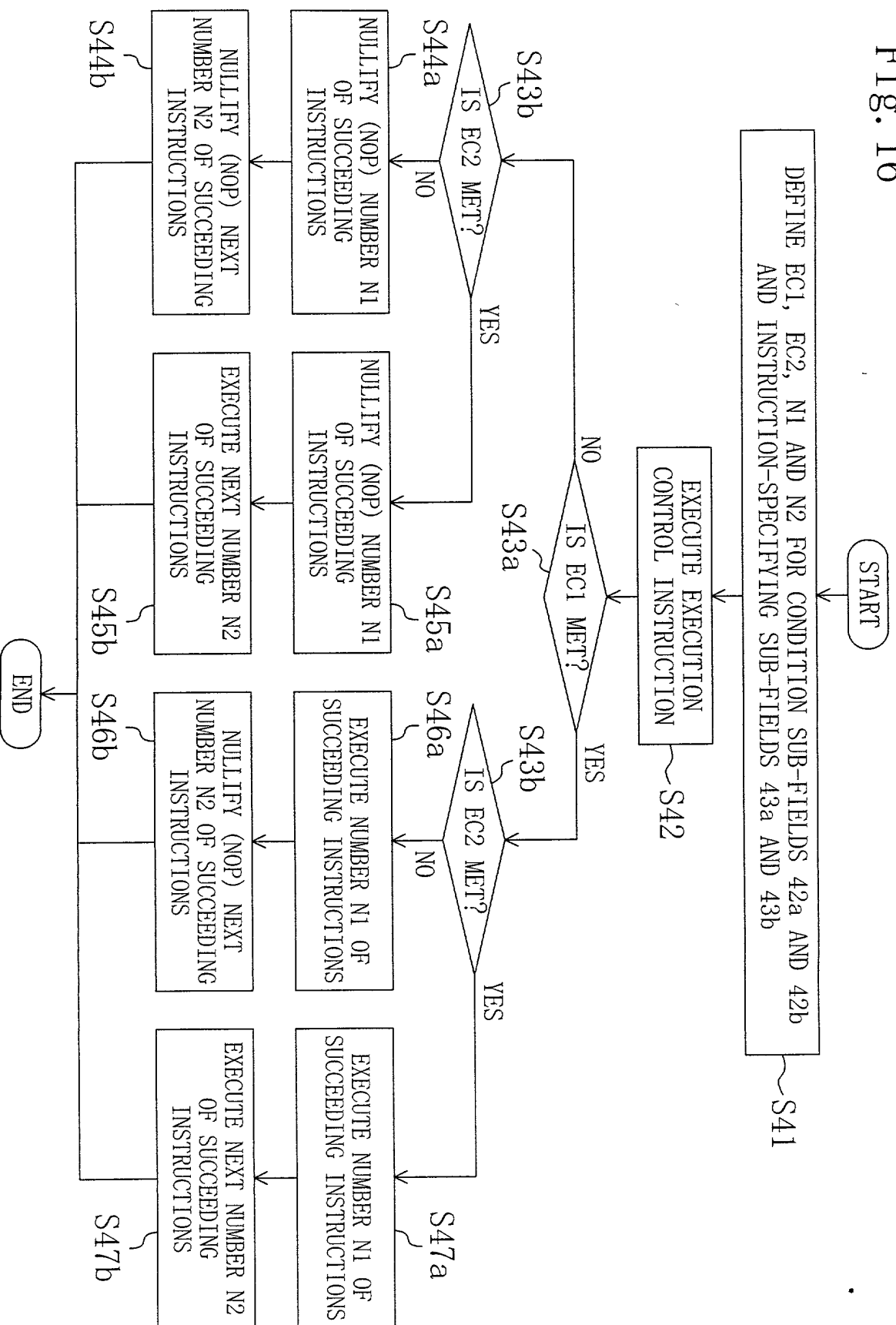


Fig. 17A

IF NEITHER EC1 NOR EC2 IS MET (WHEN $F1=0$ AND $F2=1$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 1, 3$
NOP
NOP
NOP
NOP
SUCCEEDING INSTRUCTION 5

Fig. 17B

IF EC1 IS NOT MET AND EC2 IS MET (WHEN $F1=0$ AND $F2=0$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 1, 3$
NOP
SUCCEEDING INSTRUCTION 2
SUCCEEDING INSTRUCTION 3
SUCCEEDING INSTRUCTION 4
SUCCEEDING INSTRUCTION 5

Fig. 17C

IF EC1 IS MET AND EC2 IS NOT MET (WHEN $F1=1$ AND $F2=1$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 1, 3$
SUCCEEDING INSTRUCTION 1
NOP
NOP
NOP
SUCCEEDING INSTRUCTION 5

Fig. 17D

IF BOTH EC1 AND EC2 ARE MET (WHEN $F1=1$ AND $F2=0$)
EXECUTION CONTROL INSTRUCTION $F1, NF2, 1, 3$
SUCCEEDING INSTRUCTION 1
SUCCEEDING INSTRUCTION 2
SUCCEEDING INSTRUCTION 3
SUCCEEDING INSTRUCTION 4
SUCCEEDING INSTRUCTION 5

Docket No. _____

COMBINED DECLARATION/POWER OF ATTORNEY
FOR PATENT APPLICATION

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name.

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled
SYSTEM AND METHOD FOR CONTROLLING PROGRAM EXECUTION,
the specification of which

(check one) X is attached hereto.

_____ was filed on _____ as
Application Serial No. _____.

I hereby state that I have reviewed and understand the contents of the above identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose information which is material to the examination of this application in accordance with Title 37, Code of Federal Regulations, § 1.56(a).

I hereby claim foreign priority benefits under Title 35, United States Code, § 119 of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on which priority is claimed:

Prior Foreign Application(s)			Priority Claimed
<u>11-029568</u> (Number)	<u>JAPAN</u> (Country)	<u>08/02/1999</u> (Day/Month/Year Filed)	<u> X </u> Yes <u> </u> No
_____ (Number)	_____ (Country)	_____ (Day/Month/Year Filed)	<u> </u> Yes <u> </u> No
_____ (Number)	_____ (Country)	_____ (Day/Month/Year Filed)	<u> </u> Yes <u> </u> No

I hereby claim the benefit under Title 35, United States Code, § 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, § 112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, § 1.56(a) which occurred between the filing date of the prior application and the national or PCT international filing date of this application:

(Appln. Serial No.)	(Filing Date)	(Status-patented, pending, abandoned)
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(Appln. Serial No.)	(Filing Date)	(Status-patented, pending, abandoned)
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I hereby appoint as my attorneys, with full power of substitution and revocation, to prosecute the patent application identified above and to transact all business in the U.S. Patent and Trademark Office connected therewith: Raphael V. Lupo (Reg. No. 28,363); Jack Q. Lever, Jr. (Reg. No. 28,149); Kenneth L. Cage (Reg. No. 26,151); Stanislaus Aksman (Reg. No. 28,562); Paul Devinsky (Reg. No. 28,553); Edward E. Kubasiewicz (Reg. No. 30,020), Michael E. Fogarty (Reg. No. 36,139); Brian E. Ferguson (Reg. No. 36,801); Robert W. Zelnick (Reg. No. 36,976); and Wilhlem F. Gadiano (Reg. No. 37,136).

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The undersigned hereby authorizes the U.S. attorneys named herein to accept and follow instructions from Maeda Patent Office as to any action to be taken in the Patent and Trademark Office regarding this application without direct communication between the U.S. attorney and the undersigned. In the event of a change in the persons from whom instructions may be taken, the U.S. attorneys named herein will be so notified by the undersigned.

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

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Docket No.: 43889-916

PATENT

, IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re Application of :

Masayuki YAMASAKI, et al. :

Serial No.: : Group Art Unit:

Filed: February 08, 2000 : Examiner:

For: SYSTEM AND METHOD FOR CONTROLLING PROGRAM EXECUTION

ASSOCIATE POWER OF ATTORNEY

Honorable Commissioner of
Patents and Trademarks
Washington, D. C. 20231

Sir:

The undersigned Principal Attorney of record hereby appoints the following Attorneys as his Associates with regard to the above-identified application: Edward A. Becker, Reg. No. 37,777; Stephen A. Becker, Reg. No. 26,527; Marcel K. Bingham, Reg. No. 42,327; John G. Bisbikis, Reg. No. 37,095; Carl L. Brandt, Reg. No. 44,555; Daniel Bucca, Reg. No. 42,368; Kenneth L. Cage, Reg. No. 26,151; Carina M. Tan, Reg. No. P-45,769; Stephen C. Carlson, Reg. No. 39,929; Jennifer Chen, Reg. No. 42,404; Thomas A. Corrado, Reg. No. 42,439; Lawrence T. Cullen, Reg. No. 44,489; Paul Devinsky, Reg. No. 28,553; Margaret M. Duncan, Reg. No. 30,879; Brian E. Ferguson, Reg. No. 36,801; Michael E. Fogarty, Reg. No. 36,139; John R. Fuisz, Reg. No. 37,327; Willem F. Gadiano, Reg. No. 37,136; Keith E. George, Reg. No. 34,111; Matthew V. Grumbling, Reg. No. 44,427; John A. Hankins, Reg. No. 32,029; Brian D. Hickman, Reg. No. 35,894; Joseph Hyosuk Kim, Reg. No. 41,425; Eric J. Kraus, Reg. No. 36,190; Jack Q. Lever, Reg. No. 28,149; Raphael V. Lupo, Reg. No. 28,363; Christine F. Martin, Reg. No. 39,762; Michael A. Messina, Reg. No. 33,424; Eugene J. Molinelli, Reg. No. 42,901; Christopher J. Palermo, Reg. No. 42,056; Dawn L. Palmer, Reg. No. 41,238; Joseph H. Paquin, Jr., Reg. No. 31,647; William D. Pegg, Reg. No. 42,988; Robert L. Price, Reg. No. 22,685; Thomas D. Robbins, Reg. No. 43,669; Gene Z. Robinson, Reg. No. 33,351; Joy Ann G. Serauskas, Reg. No. 27,952; David A. Spenard, Reg. No. 37,449; Arthur J. Steiner, Reg. No. 26,106; David L. Stewart, Reg. No. 37,578; Michael D. Switzer, Reg. No. 39,552; Leonid D. Thenor, Reg. No. 39,397; Keith J. Townsend, Reg. No. 40,358; Daniel S. Trainor, Reg. No. 43,959; Cameron K.

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